
Graduate Texts in Mathematics 86

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J. H. van Lint

Introduction to Coding Theory

Third Revised and Expanded Edition



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Preface to the Third Edition

It is gratifying that this textbook is still sufficiently popular to warrant a third edition. I have used the opportunity to improve and enlarge the book.

When the second edition was prepared, only two pages on algebraic geometry codes were added. These have now been removed and replaced by a relatively long chapter on this subject. Although it is still only an introduction, the chapter requires more mathematical background of the reader than the remainder of this book.

One of the very interesting recent developments concerns binary codes defined by using codes over the alphabet \mathbb{Z}_4 . There is so much interest in this area that a chapter on the essentials was added. Knowledge of this chapter will allow the reader to study recent literature on \mathbb{Z}_4 -codes.

Furthermore, some material has been added that appeared in my Springer Lecture Notes 2011, but was not included in earlier editions of this book, e. g. Generalized Reed-Solomon Codes and Generalized Reed-Muller Codes. In Chapter 2, a section on "Coding Gain" (the engineer's justification for using error-correcting codes) was added.

For the author, preparing this third edition was a most welcome return to mathematics after seven years of administration. For valuable discussions on the new material, I thank C. P. J. M. Baggen, I. M. Duursma, H. D. L. Hollmann, H. C. A. van Tilborg, and R. M. Wilson. A special word of thanks to R. A. Pellikant for his assistance with Chapter 10.

Eindhoven
November 1998

J. H. VAN LINT

Preface to the Second Edition

The first edition of this book was conceived in 1981 as an alternative to outdated, oversized, or overly specialized textbooks in this area of discrete mathematics - a field that is still growing in importance as the need for mathematicians and computer scientists in industry continues to grow.

The body of the book consists of two parts: a rigorous, mathematically oriented first course in coding theory followed by introductions to special topics. The second edition has been largely expanded and revised. The main additions in the second edition are:

- (1) a long section on the binary Golay code;
- (2) a section on Kerckhoff's codes;
- (3) a treatment of the Van Lint-Wilson bound for the minimum distance of cyclic codes;
- (4) a section on binary cyclic codes of even length;
- (5) an introduction to algebraic geometry codes.

Eindhoven
November 1991

J.H. VAN LINT

Preface to the First Edition

Coding theory is still a young subject. One can safely say that it was born in 1948. It is not surprising that it has not yet become a fixed topic in the curriculum of most universities. On the other hand, it is obvious that discrete mathematics is rapidly growing in importance. The growing need for mathematicians and computer scientists in industry will lead to an increase in courses offered in the area of discrete mathematics. One of the most suitable and fascinating is, indeed, coding theory. So, it is not surprising that one more book on this subject now appears. However, a little more justification and a little more history of the book are necessary. At a meeting on coding theory in 1979 it was remarked that there was no book available that could be used for an introductory course on coding theory (mainly for mathematicians but also for students in engineering or computer science). The best known textbooks were either too old, too big, too technical, too much for specialists, etc. The final remark was that my Springer Lecture Notes (#203) were slightly obsolete and out of print. Without realizing what I was getting into I announced that the statement was not true and proved this by showing several participants the book *Inleiding in de Coderingstheorie*, a little book based on the syllabus of a course given at the Mathematical Centre in Amsterdam in 1975 (M.C. Syllabus 31). The course, which was a great success, was given by M.R. Best, A.L. Brouwer, P. van Emde Boas, T.M.V. Janssen, H.W. Lenstra Jr., A. Schrijver, H.C.A. van Tilborg and myself. Since then the book has been used for a number of years at the Technological Universities of Delft and Eindhoven.

The comments above explain why it seemed reasonable (to me) to translate the Dutch book into English. In the name of Springer-Verlag I thank the Mathematical Centre in Amsterdam for permission to do so. Of course it turned out to be more than a translation. Much was rewritten or expanded,

problems were changed and solutions were added, and a new chapter and several new proofs were included. Nevertheless the M.C. Syllabus (and the Springer Lecture Notes 201) are the basis of this book.

The book consists of three parts. Chapter 1 contains the prerequisite mathematical knowledge. It is written in the style of a memory-refresher. The reader who discovers topics that he does not know will get some idea about them but it is recommended that he also looks at standard textbooks on those topics. Chapters 2 to 6 provide an introductory course in coding theory. Finally, Chapters 7 to 11 are introductions to special topics and can be used as supplementary reading or as a preparation for studying the literature.

Despite the youth of the subject, which is demonstrated by the fact that the papers mentioned in the references have 1974 as the average publication year, I have not considered it necessary to give credit to every author of the theorems, lemmas, etc. Some have simply become standard knowledge.

It seems appropriate to mention a number of textbooks that I use regularly and that I would like to recommend to the student who would like to learn more than this introduction can offer. First of all P.J. MacWilliams and N.J.A. Sloane, *The Theory of Error-Correcting Codes* (reference [46]), which contains a much more extensive treatment of most of what is in this book and has 1500 references! For the more technically oriented student with an interest in decoding, complexity questions, etc. E.R. Berlekamp's *Algebraic Coding Theory* (reference [2]) is a must. For a very well-written mixture of information theory and coding theory I recommend: R.J. McEliece, *The Theory of Information and Coding* (reference [51]). In the present book very little attention is paid to the relation between coding theory and combinatorial mathematics. For this the reader should consult P.J. Cameron and J.H. van Lint, *Designs, Graphs, Codes and their Links* (reference [11]).

I sincerely hope that the time spent writing this book (instead of doing research) will be considered well invested.

Eindhoven
July 1981

J.H. VAN LINT

Second edition comments: Apparently the hope expressed in the final line of the preface of the first edition came true: a second edition has become necessary. Several misprints have been corrected and also some errors. In a few places some extra material has been added.

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CHAPTER 1

Mathematical Background

In order to be able to read this book a fairly thorough mathematical background is necessary. In different chapters many different areas of mathematics play a rôle. The most important one is certainly algebra but the reader must also know some facts from elementary number theory, probability theory and a number of concepts from combinatorial theory such as designs and geometries. In the following sections we shall give a brief survey of the prerequisite knowledge. Usually proofs will be omitted. For these we refer to standard textbooks. In some of the chapters we need a large number of facts concerning a not too well-known class of orthogonal polynomials, called Krawtchouk polynomials. These properties are treated in Section 1.2. The notations that we use are fairly standard. We mention a few that may not be generally known. If C is a finite set we denote the number of elements of C by $|C|$. If the expression B is the definition of concept A then we write $A := B$. We use "iff" for "if and only if". An identity matrix is denoted by I and the matrix with all entries equal to one is J . Similarly we abbreviate the vector with all coordinates 0 (resp. 1) by $\mathbf{0}$ (resp. $\mathbf{1}$). Instead of using $\lfloor x \rfloor$ we write $\lfloor x \rfloor := \max\{n \in \mathcal{Z} : n \leq x\}$ and we use the symbol $\lceil x \rceil$ for rounding upwards.

§1.1. Algebra

We need only very little from elementary number theory. We assume known that in \mathbb{N} every number can be written in exactly one way as a product of prime numbers (if we ignore the order of the factors). If a divides b , then we write $a \mid b$. If p is a prime number and $p^r \mid a$ but $p^{r+1} \nmid a$, then we write $p^r \parallel a$. If

$k \in \mathbb{N}$, $k > 1$, then a representation of n in the base k is a representation

$$n = \sum_{i=0}^l n_i k^i,$$

$0 \leq n_i < k$ for $0 \leq i \leq l$. The largest integer π such that $\pi|a$ and $\pi|b$ is called the greatest common divisor of a and b and denoted by $\text{g.c.d.}(a, b)$ or simply (a, b) . If $m|(a - b)$ we write $a \equiv b \pmod{m}$.

(1.1.1) Theorem. *If*

$$\varphi(n) := \#\{m \in \mathbb{N} \mid 1 \leq m < n, (m, n) = 1\},$$

then

- (i) $\varphi(n) = n \prod_{p|n} (1 - 1/p)$,
- (ii) $\sum_{d|n} \varphi(d) = n$.

The function φ is called the *Euler indicator*.

(1.1.2) Theorem. *If $(a, m) = 1$ then $a^{\varphi(m)} \equiv 1 \pmod{m}$.*

Theorem 1.1.2 is called the Euler–Fermat theorem.

(1.1.3) Definition. The *Möbius function* μ is defined by

$$\mu(n) := \begin{cases} 1, & \text{if } n = 1, \\ (-1)^k, & \text{if } n \text{ is the product of } k \text{ distinct prime factors,} \\ 0, & \text{otherwise.} \end{cases}$$

(1.1.4) Theorem. *If f and g are functions defined on \mathbb{N} such that*

$$g(n) = \sum_{d|n} f(d),$$

then

$$f(n) = \sum_{d|n} \mu(d)g\left(\frac{n}{d}\right).$$

Theorem 1.1.4 is known as the *Möbius inversion formula*.

Algebraic Structures

We assume that the reader is familiar with the basic ideas and theorems of linear algebra although we do refresh his memory below. We shall first give a sequence of definitions of algebraic structures with which the reader must be familiar in order to appreciate algebraic coding theory.

(1.1.5) **Definition.** A *group* (G, \cdot) is a set G on which a product operation has been defined satisfying

- (i) $\forall a, b \in G [ab \in G]$;
- (ii) $\forall a, b, c \in G [a(bc) = (ab)c]$;
- (iii) $\exists e \in G \forall a \in G [ae = ea = a]$,
(the element e is unique);
- (iv) $\forall a \in G \exists a^{-1} \in G [aa^{-1} = a^{-1}a = e]$,
(a^{-1} is called the inverse of a and also denoted by a^{-1}).

If furthermore

- (v) $\forall a, b \in G [ab = ba]$.

then the group is called *abelian* or *commutative*.

If (G, \cdot) is a group and $H \subseteq G$ such that (H, \cdot) is also a group, then (H, \cdot) is called a *subgroup* of (G, \cdot) . Usually we write G instead of (G, \cdot) . The number of elements of a finite group is called the *order* of the group. If (G, \cdot) is a group and $a \in G$, then the smallest positive integer n such that $a^n = e$ (if such an n exists) is called the *order* of a . In this case the elements $e, a, a^2, \dots, a^{n-1}$ form a so-called *cyclic subgroup* with a as *generator*. If (G, \cdot) is abelian and (H, \cdot) is a subgroup then the sets $aH := \{ah | h \in H\}$ are called *cosets* of H . Since two cosets are obviously disjoint or identical, the cosets form a partition of G . An element chosen from a coset is called a *representative* of the coset. It is not difficult to show that the cosets again form a group if we define multiplication of cosets by $(aH)(bH) := abH$. This group is called the *factor group* and indicated by G/H . As a consequence note that if $a \in G$, then the order of a divides the order of G (also if G is not abelian).

A fundamental theorem of group theory states that a finite abelian group is a direct sum of cyclic groups.

(1.1.6) **Definition.** A set R with two operations, usually called addition and multiplication, denoted by $(R, +, \cdot)$, is called a *ring* if

- (i) $(R, +)$ is an abelian group.
- (ii) $\forall a, b \in R \forall c \in R [(ab)c = a(bc)]$.
- (iii) $\forall a, b, c \in R [a(b+c) = ab + ac \wedge (a+b)c = ac + bc]$.

The identity element of $(R, +)$ is usually denoted by 0 .

If the additional property

- (iv) $\forall a, b \in R [ab = ba]$

holds, then the ring is called *commutative*.

The integers \mathcal{Z} are the best known example of a ring.

If $(R, +, \cdot)$ is a commutative ring, a nonzero element $a \in R$ is called a *zero divisor* if there exists a nonzero element $b \in R$ such that $ab = 0$. If a nontrivial

ring has no zero divisors, it is called an *integral domain*. In the same way that \mathcal{Z} is extended to \mathbb{Q} , an integral domain can be embedded in its *field of fractions* or *quotient field*.

(1.1.7) **Definition.** If $(R, +, \cdot)$ is a ring and $\emptyset \neq S \subseteq R$, then S is called an *ideal* if

- (i) $\forall a, b \in S, a + b \in S$,
- (ii) $\forall a, b \in S, ab \in S$ and $ba \in S$.

It is clear that if S is an ideal in R , then $(S, +, \cdot)$ is a subring, but requirement (ii) says more than that.

(1.1.8) **Definition.** A *field* is a ring $(R, +, \cdot)$ for which $(R \setminus \{0\}, \cdot)$ is an abelian group.

(1.1.9) **Theorem.** Every finite ring R with at least two elements such that

$$\forall a, b \in R, ab = 0 \Rightarrow (a = 0 \vee b = 0)$$

is a field.

(1.1.10) **Definition.** Let $(V, +)$ be an abelian group, F a field and let a multiplication $V \times V \rightarrow V$ be defined satisfying

- (i) $\forall a \in V, [1a = a]$,
 $\forall \alpha, \beta \in F, \forall a \in V, \alpha(\beta a) = (\alpha\beta)a$,
- (ii) $\forall \alpha, \beta \in F, \forall a, b \in V, \alpha(a + b) = \alpha a + \alpha b$,
 $\forall \alpha \in F, \forall a, b \in V, (\alpha + \beta)a = \alpha a + \beta a$.

Then the triple $(V, +, F)$ is called a *vector space* over the field F . The identity element of $(V, +)$ is denoted by $\mathbf{0}$.

We assume the reader to be familiar with the vector space \mathbb{R}^n consisting of all n -tuples (a_1, a_2, \dots, a_n) with the obvious rules for addition and multiplication. We remind him of the fact that a *k-dimensional subspace* C of this vector space is a vector space with a basis consisting of vectors $\mathbf{a}_1 := (a_{11}, a_{12}, \dots, a_{1n}), \mathbf{a}_2 := (a_{21}, a_{22}, \dots, a_{2n}), \dots, \mathbf{a}_k := (a_{k1}, a_{k2}, \dots, a_{kn})$, where the word *basis* means that every $\mathbf{a} \in C$ can be written in a unique way as $x_1\mathbf{a}_1 + x_2\mathbf{a}_2 + \dots + x_k\mathbf{a}_k$. The reader should also be familiar with the process of going from one basis of C to another by taking combinations of basis vectors, etc. We shall usually write vectors as row vectors as we did above. The *inner product* $\langle \mathbf{a}, \mathbf{b} \rangle$ of two vectors \mathbf{a} and \mathbf{b} is defined by

$$\langle \mathbf{a}, \mathbf{b} \rangle := a_1b_1 + a_2b_2 + \dots + a_nb_n.$$

The elements of a basis are called *linearly independent*. In other words this means that a linear combination of these vectors is $\mathbf{0}$ iff all the coefficients are 0. If $\mathbf{a}_1, \dots, \mathbf{a}_k$ are k linearly independent vectors, i.e. a basis of a k -dimensional

subspace C , then the system of equations $(a_i, y) = 0$ ($i = 1, 2, \dots, k$) has as its solution all the vectors in a subspace of dimension $n - k$ which we denote by C' . So,

$$C' := \{y \in \mathbb{R}^n \mid \forall_{a_i \in C} \langle a_i, y \rangle = 0\}.$$

These ideas play a fundamental role later on, where \mathbb{R} is replaced by a finite field F . The theory reviewed above goes through in that case.

(1.1.11) Definition. Let $(V, +)$ be a vector space over \mathbb{F} and let α multiplication $V \times V \rightarrow V$ be defined that satisfies

- (i) $(V, +, \alpha)$ is a ring.
- (ii) $\forall_{a_i \in C} \forall_{x \in V} \forall_{y \in V} [(xa)y = \alpha(xy)]$.

Then we say that the system is an *algebra* over F .

Suppose we have a finite group (G, \cdot) and we consider the elements of G as basis vectors for a vector space $(V, +)$ over a field F . Then the elements of V are represented by linear combinations $x_1g_1 + x_2g_2 + \dots + x_n g_n$, where

$$x_i \in F, \quad g_i \in G, \quad (1 \leq i \leq n = |G|).$$

We can define a multiplication \bullet for these vectors in the obvious way, namely

$$\left(\sum_i x_i g_i\right) \bullet \left(\sum_j y_j g_j\right) := \sum_i \sum_j (x_i y_j) (g_i \cdot g_j),$$

which can be written as $\sum_k y_k g_k$, where y_k is the sum of the elements $x_i y_j$ over all pairs (i, j) such that $g_i \cdot g_j = g_k$. This yields an algebra which is called the *group algebra* of G over F and denoted by FG .

EXAMPLES. Let us consider a number of examples of the concepts defined above.

If $A := \{a_1, a_2, \dots, a_n\}$ is a finite set, we can consider all one-to-one mappings of S onto S . These are called *permutations*. If σ_1 and σ_2 are permutations we define $\sigma_1 \sigma_2$ by $(\sigma_1 \sigma_2)(a) := \sigma_1(\sigma_2(a))$ for all $a \in A$. It is easy to see that the set S_n of all permutations of A with this multiplication is a group, known as the *symmetric group of degree n* . In this book we shall often be interested in special permutation groups. These are subgroups of S_n . We give one example. Let C be a k -dimensional subspace of \mathbb{R}^n . Consider all permutations σ of the integers $1, 2, \dots, n$ such that for every vector $v = (c_1, c_2, \dots, c_n) \in C$ the vector $(c_{\sigma(1)}, c_{\sigma(2)}, \dots, c_{\sigma(n)})$ is also in C . These clearly form a subgroup of S_n . Of course C will often be such that this subgroup of S consists of the identity only but there are more interesting examples! Another example of a permutation group which will turn up later is the *affine permutation group* defined as follows. Let \mathbb{F} be a (finite) field. The mapping $f_{u,v}$, when $u \in \mathbb{F}$, $v \in \mathbb{F}$, $u \neq 0$, is defined on \mathbb{F} by $f_{u,v}(x) := ux + v$ for all $x \in \mathbb{F}$. These mappings are permutations of \mathbb{F} and clearly they form a group under composition of functions.

A *permutation matrix* P is a $\{0, 1\}$ -matrix that has exactly one 1 in each row and column. We say that P corresponds to the permutation σ of $\{1, 2, \dots, n\}$ if $p_{ij} = 1$ iff $i = \sigma(j)$ ($i = 1, 2, \dots, n$). With this convention the product of permutations corresponds to the product of their matrices. In this way one obtains the so-called matrix representation of a group of permutations.

A group G of permutations acting on a set Ω is called *k-transitive* on Ω if for every ordered k -tuple (a_1, \dots, a_k) of distinct elements of Ω and for every k -tuple (b_1, \dots, b_k) of distinct elements of Ω , there is an element $\sigma \in G$ such that $b_i = \sigma(a_i)$ for $1 \leq i \leq k$. If $k = 1$ we call the group *transitive*.

Let S be an ideal in the ring $(R, +, \cdot)$. Since $(S, +)$ is a subgroup of the abelian group $(R, +)$, we can form the factor group. The cosets are now called *residue classes mod S*. For these classes we introduce a multiplication in the obvious way: $(a + S)(b + S) := ab + S$. The reader who is not familiar with this concept should check that this definition makes sense (i.e. it does not depend on the choice of representatives a resp. b). In this way we have constructed a ring, called the *residue class ring* $R \bmod S$ and denoted by R/S . The following example will surely be familiar. Let $R := \mathbb{Z}$ and let p be a prime. Let S be $p\mathbb{Z}$, the set of all multiples of p , which is sometimes also denoted by (p) . Then R/S is the ring of integers mod p . The elements of R/S can be represented by $0, 1, \dots, p-1$ and then addition and multiplication are the usual operations in \mathbb{Z} followed by a reduction mod p . For example, if we take $p = 7$, then $4 + 5 = 2$ because in \mathbb{Z} we have $4 + 5 = 9 \equiv 2 \pmod{7}$. In the same way $4 \cdot 5 = 6$ in $\mathbb{Z}/7\mathbb{Z} = \mathbb{Z}/(7)$. If S is an ideal in \mathbb{Z} and $S \neq \{0\}$, then there is a smallest positive integer k in S . Let $s \in S$. We can write s as $ak + b$, where $0 \leq b < k$. By the definition of ideal we have $ak \in S$ and hence $b = s - ak \in S$ and then the definition of k implies that $b = 0$. Therefore $S = \{k\}$. An ideal consisting of all multiples of a fixed element is called a *principal ideal*. If a ring R has no other ideals than principal ideals, it is called a *principal ideal ring*. Therefore \mathbb{Z} is such a ring.

An ideal S is called a *prime ideal* if $ab \in S$ implies $a \in S$ or $b \in S$. An ideal S in a ring R is called *maximal* if for every ideal I with $S \subset I \subset R$, $I = S$ or $I = R$ ($S \neq R$). If a ring has a unique maximal ideal, it is called a *local ring*.

(1.1.12) Theorem. *If p is a prime then $\mathbb{Z}/p\mathbb{Z}$ is a field.*

This is an immediate consequence of Theorem 1.1.9 but also obvious directly. A finite field with n elements is denoted by \mathbb{F}_n or $\text{GF}(n)$ (Galois field).

Rings and Finite Fields

More about finite fields will follow below. First some more about rings and ideals. Let \mathbb{F} be a finite field. Consider the set $\mathbb{F}[x]$ consisting of all polynomials $a_0 + a_1x + \dots + a_nx^n$, where n can be any integer in \mathbb{N} and $a_i \in \mathbb{F}$ for $0 \leq i \leq n$. With the usual definition of addition and multiplication of polyno-

mials this yields a ring $(F[x], +, \cdot)$ which is usually just denoted by $F[x]$. The set of all polynomials that are multiples of a fixed polynomial $g(x)$, i.e. all polynomials of the form $a(x)g(x)$ where $a(x) \in F[x]$, is an ideal in $F[x]$.

As before, we denote this ideal by $(g(x))$. The following theorem states that there are no other types.

(1.1.13) **Theorem.** $F[x]$ is a principal ideal ring.

The residue class ring $F[x]/(g(x))$ can be represented by the polynomials whose degree is less than the degree of $g(x)$. In the same way as our example \mathbb{Z}_7/\mathbb{Z} given above, we now multiply and add these representatives in the usual way and then reduce mod $g(x)$. For example, we take $\mathbb{F}_2 = \{0, 1\}$ and $g(x) = x^3 + x + 1$. Then $(x+1)(x^2+1) = x^3 + x^2 + x + 1 = x^2$. This example is a useful one to study carefully if one is not familiar with finite fields. First observe that $g(x)$ is *irreducible*, i.e., there do not exist polynomials $a(x)$ and $b(x) \in F[x]$, both of degree less than 3, such that $g(x) = a(x)b(x)$. Next, realize that this means that in $F_2[x]/(g(x))$ the product of two elements $a(x)$ and $b(x)$ is 0 iff $a(x) = 0$ or $b(x) = 0$. By Theorem 1.1.9 this means that $F_2[x]/(g(x))$ is a field. Since the representatives of this residue class ring all have degrees less than 3, there are exactly eight of them. So we have found a field with eight elements, i.e. F_{2^3} . This is an example of the way in which finite fields are constructed.

(1.1.14) **Theorem.** Let p be a prime and let $g(x)$ be an irreducible polynomial of degree r in the ring $F_p[x]$. Then the residue class ring $F_p[x]/(g(x))$ is a field with p^r elements.

PROOF. The proof is the same as the one given for the example $p = 2, r = 3, g(x) = x^3 + x + 1$. □

(1.1.15) **Theorem.** Let F be a field with n elements. Then n is a power of a prime.

PROOF. By definition there is an identity element for multiplication in F . We denote this by 1. Of course $1 + 1 \in F$ and we denote this element by 2. We continue in this way, i.e. $2 + 1 = 3$, etc. After a finite number of steps we encounter a field element that already has a name. Suppose, e.g. that the sum of k terms 1 is equal to the sum of l terms 1 ($k > l$). Then the sum of $(k - l)$ terms 1 is 0, i.e. the first time we encounter an element that already has a name, this element is 0. Say 0 is the sum of k terms 1. If k is composite, $k = ab$, then the product of the elements which we have called a resp. b is 0, a contradiction. So k is a prime and we have shown that F_p is a subfield of F . We define linear independence of a set of elements of F with respect to (coefficients from) F_p in the obvious way. Among all linearly independent subsets of F let $\{x_1, x_2, \dots, x_r\}$ be one with the maximal number of elements. If x is any element of F then the elements x, x_1, x_2, \dots, x_r are not linearly

independent, i.e. there are coefficients $0 \neq a_1, \dots, a_r$ such that $a_1x_1 + \dots + a_r x_r = 0$ and hence x is a linear combination of x_1 to x_r . Since there are obviously p^r distinct linear combinations of x_1 to x_r , the proof is complete. \square

From the previous theorems we now know that a field with n elements exists iff n is a prime power, providing we can show that for every $r \geq 1$ there is an irreducible polynomial of degree r in $\mathbb{F}_p[x]$. We shall prove this by calculating the number of such polynomials (fix p and let I_r denote the number of irreducible polynomials of degree r that are *monic*, i.e. the coefficient of x^r is 1. We claim that

$$(1.1.16) \quad (1 - px)^{-1} = \prod_{r=1}^{\infty} (1 - x^r)^{-I_r}.$$

In order to see this, first observe that the coefficient of x^n on the left-hand side is p^n , which is the number of monic polynomials of degree n with coefficients in \mathbb{F}_p . We know that each such polynomial can be factored uniquely into irreducible factors and we must therefore convince ourselves that these products are counted on the right-hand side of (1.1.16). To show this we just consider two irreducible polynomials $a_1(x)$ of degree r and $a_2(x)$ of degree s . There is a 1-1 correspondence between products $(a_1(x))^i (a_2(x))^j$ and terms $x_1^i x_2^j$ in the product of $(1 + x_1 + x_1^2 + \dots)$ and $(1 + x_2 + x_2^2 + \dots)$. If we identify x_1 and x_2 with x , then the exponent of x is the degree of $(a_1(x))^i (a_2(x))^j$. Instead of two polynomials $a_1(x)$ and $a_2(x)$, we now consider all irreducible polynomials and (1.1.16) follows.

In (1.1.16) we take logarithms on both sides, then differentiate, and finally multiply by x to obtain

$$(1.1.17) \quad \frac{px}{1 - px} = \sum_{r=1}^{\infty} I_r \frac{x^r}{1 - x^r}.$$

Comparing coefficients of x^d on both sides of (1.1.17) we find

$$(1.1.18) \quad p^d = \sum_{r|d} r I_r.$$

Now apply Theorem 1.1.4 to (1.1.18). We find

$$(1.1.19) \quad I_r = \frac{1}{r} \sum_{d|r} \mu(d) p^{d/r} \geq \frac{1}{r} (p^{r/r} - p^{r/2} - p^{r/3} - \dots) \\ \geq \frac{1}{r} \left(p^r - \sum_{\substack{d|r \\ d < r}} p^d \right) \geq \frac{1}{r} p^r (1 - p^{-r/2} - 1) \geq 0.$$

Now that we know for which values of n a field with n elements exists, we wish to know more about these fields. The structure of \mathbb{F}_p will play a very important role in many chapters of this book. As a preparation consider a finite field \mathbb{F} and a polynomial $f(x) \in \mathbb{F}[x]$ such that $f(a) = 0$, where $a \in \mathbb{F}$. Then by dividing we find that there is a $g(x) \in \mathbb{F}[x]$ such that $f(x) = (x - a)g(x)$.

Continuing in this way we establish the trivial fact that a polynomial $f(x)$ of degree r in $\mathbb{F}_q[x]$ has at most r zeros in \mathbb{F}_q .

If α is an element of order e in the multiplicative group $(\mathbb{F}_q \setminus \{0\}, \cdot)$, then α is a zero of the polynomial $x^e - 1$. In fact, we have

$$x^e - 1 = (x - 1)(x - \alpha)(x - \alpha^2) \cdots (x - \alpha^{e-1}).$$

It follows that the only elements of order e in the group are the powers α^i where $1 \leq i < e$ and $(i, e) = 1$. There are $\phi(e)$ such elements. Hence, for every e which divides $p^r - 1$ there are either 0 or $\phi(e)$ elements of order e in the field. By (1.1.1) the possibility 0 never occurs. As a consequence there are elements of order $p^r - 1$, in fact exactly $\phi(p^r - 1)$ such elements. We have proved the following theorem.

(1.1.20) Theorem. *In \mathbb{F}_q the multiplicative group $(\mathbb{F}_q \setminus \{0\}, \cdot)$ is a cyclic group.*

This group is often denoted by \mathbb{F}_q^* .

(1.1.21) Definition. A generator of the multiplicative group of \mathbb{F}_q is called a *primitive element* of the field.

Note that Theorem 1.1.20 states that the elements of \mathbb{F}_q are exactly the q distinct zeros of the polynomial $x^q - x$. An element β such that $\beta^k = 1$ but $\beta^l \neq 1$ for $0 < l < k$ is called a *primitive k th root of unity*. Clearly a primitive element α of \mathbb{F}_q is a primitive $(q - 1)$ th root of unity. If e divides $q - 1$ then α^e is a primitive $((q - 1)/e)$ th root of unity. Furthermore a consequence of Theorem 1.1.20 is that \mathbb{F}_r is a subfield of \mathbb{F}_q iff r divides q . Actually this statement could be slightly confusing to the reader. We have been suggesting by our notation that for a given q the field \mathbb{F}_q is unique. This is indeed true. In fact this follows from (1.1.18). We have shown that for $q = p^n$ every element of \mathbb{F}_q is a zero of some irreducible factor of $x^q - x$ and from the remark above and Theorem 1.1.14 we see that this factor must have a degree r such that $r|n$. By (1.1.18) this means we have used all irreducible polynomials of degree r where $r|n$. In other words, the product of these polynomials is $x^q - x$. This establishes the fact that two fields \mathbb{F} and \mathbb{F}' of order q are isomorphic, i.e. there is a mapping $\varphi: \mathbb{F} \rightarrow \mathbb{F}'$ which is one-to-one and such that φ preserves addition and multiplication.

The following theorem is used very often in this book.

(1.1.22) Theorem. *Let $q = p^r$ and $0 \neq f(x) \in \mathbb{F}_q[x]$.*

- (i) *If $\alpha \in \mathbb{F}_q$ and $f(\alpha) = 0$, then $f(\alpha^q) = 0$.*
- (ii) *Conversely: Let $g(x)$ be a polynomial with coefficients in an extension field of \mathbb{F}_q . If $g(\alpha^q) = 0$ for every α for which $g(\alpha) = 0$, then $g(x) \in \mathbb{F}_q[x]$.*

PROOF.

- (i) By the binomial theorem we have $(a + b)^p = a^p + b^p$ because p divides $\binom{p}{k}$ for $1 < k < p - 1$. It follows that $(a + b)^p = a^p + b^p$. If $f(x) = \sum a_i x^i$ then $(f(x))^p = \sum a_i^p (x^i)^p$. Because $a_i \in \mathbb{F}_q$ we have $a_i^p = a_i$. Substituting $x = \alpha$ we find $f(\alpha^p) = (f(\alpha))^p = 0$.
- (ii) We already know that in a suitable extension field of \mathbb{F}_q the polynomial $g(x)$ is a product of factors $x - \alpha_i$ (all of degree 1, that is) and if $x - \alpha_i$ is one of these factors, then $x - \alpha_i^p$ is also one of them. If $g(x) = \sum_{i=0}^{e-1} a_i x^i$ then a_k is a symmetric function of the zeros α_i and hence $a_k = a_k^p$, i.e. $a_k \in \mathbb{F}_q$.

If $\alpha \in \mathbb{F}_q$, where $q = p^e$, then the *minimal polynomial* of α over \mathbb{F}_q is the irreducible polynomial $f(x) \in \mathbb{F}_q[x]$ such that $f(\alpha) = 0$. If α has order e then from Theorem 1.1.22 we know that this minimal polynomial is $\prod_{i=0}^{e-1} (x - \alpha^{p^i})$, where m is the smallest integer such that $p^m \equiv 1 \pmod{e}$.

Sometimes we shall consider a field \mathbb{F}_q with a fixed primitive element α . In that case we use $m_\alpha(x)$ to denote the minimal polynomial of α . An irreducible polynomial which is the minimal polynomial of a primitive element in the corresponding field is called a *primitive polynomial*. Such polynomials are the most convenient ones to use in the construction of Theorem 1.1.14. We give an example in detail.

(1.1.23) EXAMPLE. The polynomial $x^4 + x + 1$ is primitive over \mathbb{F}_2 . The field \mathbb{F}_{2^4} is represented by polynomials of degree < 4 . The polynomial x is a primitive element. Since we prefer to use the symbol x for other purposes, we call this primitive element α . Note that $\alpha^4 + \alpha + 1 = 0$. Every element in \mathbb{F}_{2^4} is a linear combination of the elements $1, \alpha, \alpha^2$, and α^3 . We got the following table for \mathbb{F}_{2^4} . The reader should observe that this is the equivalent of a table of logarithms for the case of the field \mathbb{B} .

The representation on the right demonstrates again that \mathbb{F}_{2^4} can be interpreted as the vector space $(\mathbb{F}_2)^4$, where $\{1, \alpha, \alpha^2, \alpha^3\}$ is the basis. The left-hand column is easiest for multiplication (add exponents, mod 15) and the right-hand column for addition (add vectors). It is now easy to check that

$$\begin{aligned} \pi_2(x) &= (x - \alpha)(x - \alpha^2)(x - \alpha^4)(x - \alpha^8) &= x^4 + x + 1, \\ \pi_3(x) &= (x - \alpha^3)(x - \alpha^6)(x - \alpha^{12})(x - \alpha^9) &= x^4 - x^3 + x^2 + x + 1, \\ \pi_5(x) &= (x - \alpha^5)(x - \alpha^{10}) &= x^2 - x + 1, \\ \pi_7(x) &= (x - \alpha^7)(x - \alpha^{14})(x - \alpha^{13})(x - \alpha^{11}) &= x^4 - x^3 + 1, \end{aligned}$$

and the decomposition of $x^{16} - x$ into irreducible factors is

$$\begin{aligned} x^{10} - x &= x(x-1)(x^2+x-1)(x^4+x+1) \\ &\quad \times (x^4-x^2+1)(x^4+x^3+x^2+x+1). \end{aligned}$$

Note that $x^4 - x = x(x-1)(x^2+x+1)$ corresponding to the elements 0, 1, x^2 , x^{10} which form the subfield $F_4 = F_2[x]/(x^2+x+1)$. The polynomial $m_3(x)$ is irreducible but not primitive.

Table of F_2 .

0	=		=	(0 0 0 0)
1	=	1	=	(1 0 0 0)
α	=	α	=	(0 1 0 0)
α^2	=	α^2	=	(0 0 1 0)
α^3	=	α^3	=	(0 0 0 1)
α^4	=	$1 + \alpha$	=	(1 1 0 0)
α^5	=	$\alpha + \alpha^2$	=	(0 1 1 0)
α^6	=	$\alpha^2 + \alpha^3$	=	(0 0 1 1)
α^7	=	$1 + \alpha + \alpha^2$	=	(1 1 0 1)
α^8	=	$1 + \alpha^2$	=	(1 0 1 0)
α^9	=	$\alpha + \alpha^3$	=	(0 1 0 1)
α^{10}	=	$1 + \alpha + \alpha^3$	=	(1 1 1 0)
α^{11}	=	$\alpha + \alpha^2 + \alpha^3$	=	(0 1 1 1)
α^{12}	=	$1 + \alpha + \alpha^2 + \alpha^3$	=	(1 1 1 1)
α^{13}	=	$1 + \alpha^2 + \alpha^3$	=	(1 0 1 1)
α^{14}	=	$1 + \alpha^3$	=	(1 0 0 1)

The reader who is not familiar with finite fields should study (1.1.14) to (1.1.23) thoroughly and construct several examples such as F_9 , F_{27} , F_{64} with the corresponding minimal polynomials, subfields, etc. For tables of finite fields see references [9] and [10].

Polynomials

We need a few more facts about polynomials. If $f(x) \in F_q[x]$ we can define the *derivative* $f'(x)$ in a purely formal way by

$$\left(\sum_{k=0}^n a_k x^k \right)' := \sum_{k=1}^n k a_k x^{k-1}.$$

The usual rules for differentiation of sums and products go through and one finds for instance that the derivative of $(x-\alpha)^2 f(x)$ is $2(x-\alpha)f'(x) + (x-\alpha)^2 f''(x)$. Therefore the following theorem is obvious.

(1.1.24) Theorem. *If $f(x) \in F_q[x]$ and α is a multiple zero of $f(x)$ in some extension field of F_q , then α is also a zero of the derivative $f'(x)$.*

Note however, that if $q = \infty$, then the second derivative of any polynomial in $F_q[x]$ is identically 0. This tells us nothing about the multiplicity of zeros

of the polynomial). In order to get complete analogy with the theory of polynomials over \mathbb{Q} , we introduce the so-called *Hesse derivative* of a polynomial $f(x) \in \mathbb{F}_q[x]$ by

$$f^{(k)}(x) := \frac{1}{k!} f^{(k)}(x);$$

(so the k -th Hesse derivative of x^n is $\binom{n}{k} x^{n-k}$).

The reader should have no difficulty proving that x is a zero of $f(x)$ with multiplicity k iff it is a zero of $f^{(i)}(x)$ for $0 \leq i < k$ and not a zero of $f^{(k)}(x)$.

Another result to be used later is the fact that if $f(x) = \prod_{i=1}^n (x - \alpha_i)$ then $f'(x) = \sum_{i=1}^n f(x)/(x - \alpha_i)$.

The following theorem is well known.

(1.1.25) Theorem. *If the polynomials $a(x)$ and $b(x)$ in $\mathbb{F}[x]$ have greatest common divisor 1, then there are polynomials $p(x)$ and $q(x)$ in $\mathbb{F}[x]$ such that*

$$a(x)p(x) + b(x)q(x) = 1.$$

PROOF. This is an immediate consequence of Theorem 1.1.23. □

Although we know from (1.1.19) that irreducible polynomials of any degree r exist, it sometimes takes a lot of work to find one. The proof of (1.1.19) shows one way to do it. One starts with all possible polynomials of degree 1 and forms all reducible polynomials of degree 2. Any polynomial of degree 2 not in the list is irreducible. Then one proceeds in the obvious way to produce irreducible polynomials of degree 3, etc. In Section 9.2 we shall need irreducible polynomials over \mathbb{F}_2 of arbitrarily high degree. The procedure sketched above is not satisfactory for that purpose. Instead, we proceed as follows.

(1.1.26) Lemma.

$$3^{\beta+1} \mid (2^{2^\beta} + 1).$$

PROOF.

(i) For $\beta = 0$ and $f = 1$ the assertion is true.

(ii) Suppose $3^t \mid (2^{2^\beta} + 1)$. Then from

$$(2^{2^{\beta+1}} + 1) = (2^{2^\beta} + 1)((2^{2^\beta} + 1)(2^{2^\beta} - 2i + 3),$$

it follows that if $t \geq 2$, then $3^{t-1} \mid (2^{2^{\beta+1}} + 1)$. □

(1.1.27) Lemma. *If m is the order of 2 (mod 3^t), then*

$$m \cdot \varphi(3^t) = 2 \cdot 2^{t-1}.$$

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